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Motivation

- Emerging renewable energy resources
- Grids to be relieved by tariffs for the current supply
- Smart Metering helps the customer
 - Shifting tasks to times with cheaper tariffs
 - Identifying heavy electricity consumers
- Meter data needs to be transmitted, but
 - ◆ Bad connectivity for GSM/UMTS in cellars
 - Powerline communication is not available in all areas
- Wireless mesh network
- Hop-to-Hop data transport



Motivation

- Pro:
 - Wireless sensor nodes are cheap
- Contra:
 - Low Bandwidth
 - Bottleneck near Sink
- Reducing network load allows more nodes in network
- Data formats are regulated and mandatory
- ⇒ Lossless data compression is needed

Smart Meter Data

ASCII-coded (EN-62056-21)

```
F.F(00000000)
 2 0.0.0(00617827)
 3 0.0.1(00000000)
 4 0.9.2(1070813)
 5 1.8.1(00000247*kWh)
 6 1.8.1 * 04(00000172 * kWh)
 7 1.8.1 * 03(00000103 * kWh)
 8 C.51.6 * 03(00000000000000)
 9 C.51.7(00000000)
10 C.77.2(0A78422F69654D2D4952481FBF16)
11 C.86.0(20001010)
12 31.25(0.008*A)
13 51.25(0.009*A)
14 71.25(0.045*A)
15 32.25(0.66*V)
16 52.25(0.57*V)
17 72.25(225.04 * V)
18 1.35.0.01(0.000*kW)
19 P.01(0080213003000)(00000000)(15)(1)(1.5)(kW)
20
   (00.00)
   (00.02)
21
   (00.01)
23
   (02.00)
24
   (00.01)
```

Motivation

- Trade-off: Compression gain ⇔ Hardware requirements
- Sensor nodes have limited hardware resources (RAM, ROM, speed)
- Modern algorithms concentrate on compression gain

Question:

Which compression algorithm gives best results for the recorded data by using as few resources as possible?

What we did:

- Improve compression algorithms for wireless sensor nodes
- Analyze the strengths and weaknesses

Agenda

- Introduction
- Compression Algorithms
- 3 Evaluation
- 4 Conclusion

Compression Types

Entropy Coding:

▶ More

Use variable Bit-length codes for symbols, depending on their probabilities. e.g.: Huffman-, Arithmetic-, Range-Encoding

■ Wordbook Compression:

> More

Recognize repetitions of strings and use references to encode them. e.g.: LZ77, LZ78

Block-sorting:

> More

Rearrange the symbols in a way that they are easier to compress afterwards. e.g.: Burrows-Wheeler-Transformation, Move-To-Front

Reference Compression Algorithms

- Lempel Ziv Storer Szymanski (LZSS)
 Dictionary compression
- Lempel Ziv Welch (LZW)
 Dictionary compression (patent expired, used in GIF)
- DEFLATE (ZLIB, GZIP)
 Dictionary compression with Huffman Encoding (used in HTTP, PDF, ...)
- BZIP2

 Block sorting with Range-encoding
- Lempel Ziv Markov Chain Algorithm (LZMA)
 Dictionary compression with Reference History and subsequent Range-encoding

Entropy-based Compression

Adaptive Trimmed Huffman Coding (ATH)

> More

- Developed for energy-constrained wireless sensor nodes
- Adaptive Entropy Coding Scheme
- Huffman tree is trimmed to reduce memory consumption
- Prefix determines the encoding of next symbol

Adaptive Markov Chain Huffman Coding (AMCH)

> More

- Uses probabilities of successive symbols
- Each symbol has Huffman tree of following symbols
- Trees are built during compression
- Each tree has Escape symbol for not yet encountered (NYE) symbols

Wordbook-based Compression

tiny Lempel Ziv Markov Chain Algorithm (tLZMA)

- Adaption of the LZMA Scheme
- History window constrained to 128 Byte
- No Range-encoding step

Lempel Ziv Markov Chain Huffman Coding (LZMH)

- Combination of tLZMA and ATH
- Dictionary compression with 128 Byte History
- Compression of symbols with ATH method

Evaluation Methodology

Data

- 3500 ASCII-coded (EN-62056-21) datasets
 - From real smart meter installation
 - Size range from 76 to 3 100 Byte
- 95 power consumption measures of household devices
 - Sampled about once per second
 - Data in binary format
 - Sizes between 18.8 and 171.0 KByte

Methodology and Metrics

- All Methods implemented in C, no use of heap memory ZLIB, BZ2 and LZMA implementations use heap memory
- Compression rate benchmarked on a desktop PC
- Compression rate = 1 compressed data size uncompressed data size
- Processing time benchmarked on ATmega 1281 (8 kB RAM, 7.37 MHz)
- Processing time over size = consumed time uncompressed data size

Memory Consumption

Algorithm	ROM (Byte)	RAM (Byte)		
		static	stack	
LZSS	544	129	19	
LZW	550	12416	16	
ZLIB	27 960	2 690	ca. > 1 000*	
BZ2	28 332	1 564	ca. > 100 000*	
LZMA	34 442	110	ca. $> 6000000^*$	
ATH	592	170	15	
AMCH	1 680	1 820	21	
tLZMA	992	133	27	
LZMH	1 428	378	29	

^(*) Heap Memory

Memory Consumption

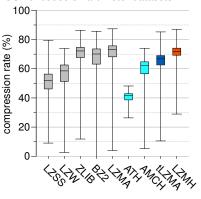
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ZLIB, BZ2, LZMA and LZW use too much memory

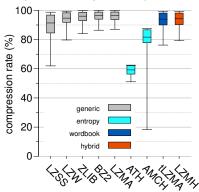
⇒ can not be tested on ATmega 1281

Compression Rate





binary encoded daily device reports

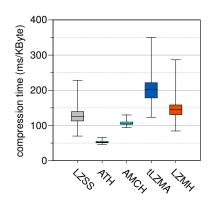


- LZMH achieves compression rates comparable to ZLIB, BZ2 and LZMA
- ATH performs good for small, but poor for bigger datasets

▶ Histogram

Processing Time

- ATH is the fastest
- ATH and AMCH are mostly unaffected by the dataset
- Other methods depend on the compressability of the datasets



▶ Histogram

Selection Guidelines

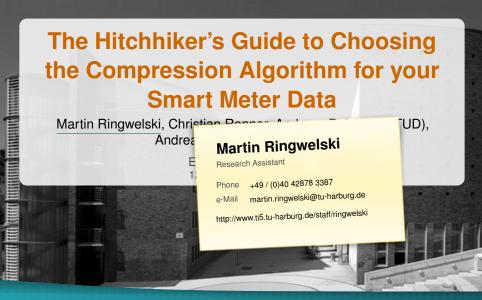
	LZSS	ATH	AMCH	tLZMA	LZMH
ROM	++	++	0	+	
RAM	++	$\oplus \oplus$	$\Theta\Theta$	++	+
Performance		$\oplus \oplus$	+	$\overline{}$	
Compression Gain	Θ	$\ominus\ominus$	0	+	+ +

Scenarios:

- Very limited resources
 - Focus on performance → ATH
 - Focus on compression → LZSS
- Limited resources, focus on compression → tLZMA
- Moderate resources, best compression → LZMH
- Good compression with predictable performance → AMCH

Resume

- Lossless data compression to reduce:
 - bandwidth
 - energy consumption
 - transmission costs
- Trade-off between compression gains, time, and resources
- ATH, LZMH, LZSS and tLZMH fulfill the resource constraints
- LZMH gives best compression results with moderate execution times and resources
- ATH is the fastest method, but has lowest compression rates



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Appendix

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Entropy

- Measure for the average self-information and the information density of a code system
- *n*: Number of symbols, p_i : Probability of a symbol $H = -\sum_{i=1}^{n} p_i \cdot \log_2 p_i$
- The entropy is the average number of needed bits to encode one symbol in a message



Move-To-Front Example

Symbol	Sequence	Alphabet
b	1	abcdefghijklmnopqrstuvwxyz
а	1	bacdefghijklmnopqrstuvwxyz
n	13	abcdefghijklmnopqrstuvwxyz
а	1	nabcdefghijklmopqrstuvwxyz
n	1	anbcdefghijklmopqrstuvwxyz
a	1	nabcdefghijklmopqrstuvwxyz

▶ Back

LZSS Example

```
Hello you there, yea you there, yea you there, Hello!
```

Becomes:

```
Hello you there, yea(5,11)(18,15)(0,5)!
```

The Numbers in brackets determine the offset and length of the previous occurrence of the following string.

▶ Back

Range-encoding Example

Message: AABA\$

Intervals: A: 0 to 0,6 - B: 0,6 to 0,8 - \$: 0,8 to 1

A - 0 to 0,6

A - 0 to 0,36

B - 0,216 to 0,288

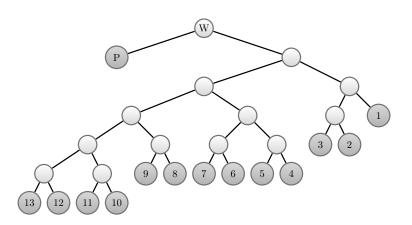
A - 0,216 to 0,2591

\$ - 0,25046 to 0,2591

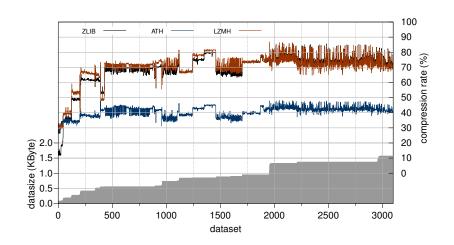
 \Rightarrow 0,251 is the shortest number that lies in the end interval and thereby encodes the message AABA\$.

▶ Back

The used ATH tree

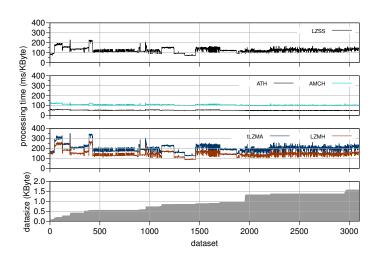


Histogram



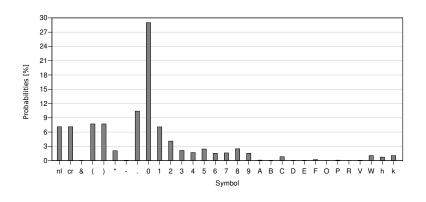


Histogram of Times





Symbol Histogram



Symbol-transition Heatmap

